

Applied Artificial Intelligence

Session 4: Searching as an AI Method

Fall 2018

NC State University

Instructor: Dr. Behnam Kia

Course Website: <https://appliedai.wordpress.ncsu.edu/>

Computational Complexity & Complexity Classes

Computational Complexity

Computational complexity theory focuses on classifying computational problems according to their inherent difficulty, and relating the resulting complexity classes to each other.

– Wikipedia

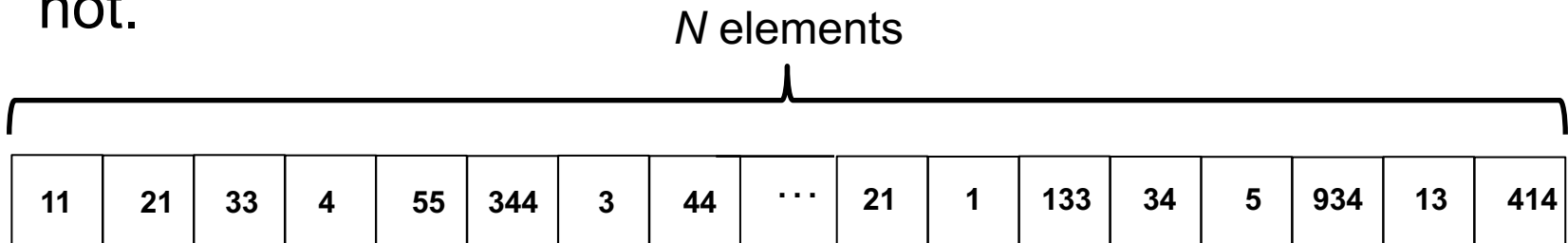
Computational Complexity:

Why We Need It?

- We would like to get an idea of hardness of a problem we like to solve.
- It tells us how many resources we need to solve the problem.
- Or is it even practical trying to directly solve it?

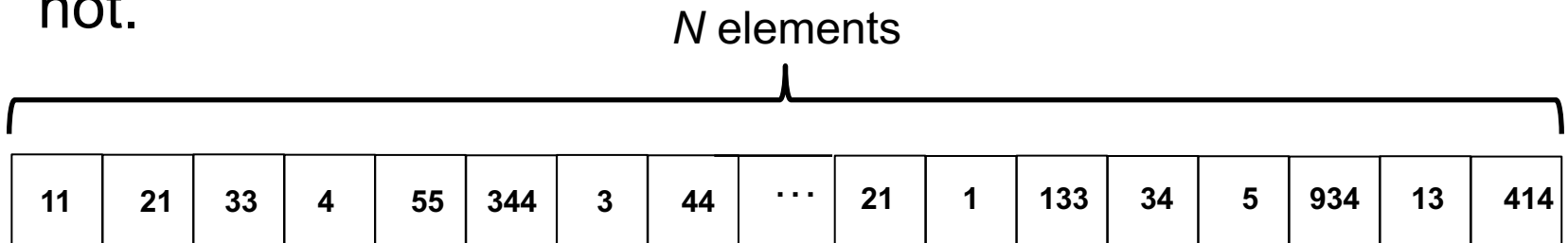
Example: Searching an Array of Size N for a specific Element

- An array of size n is filled with unsorted randomly placed, numbers.
- Goal: Find whether a specific item x is in this array or not.



Searching an Array of Size N for A specific Element

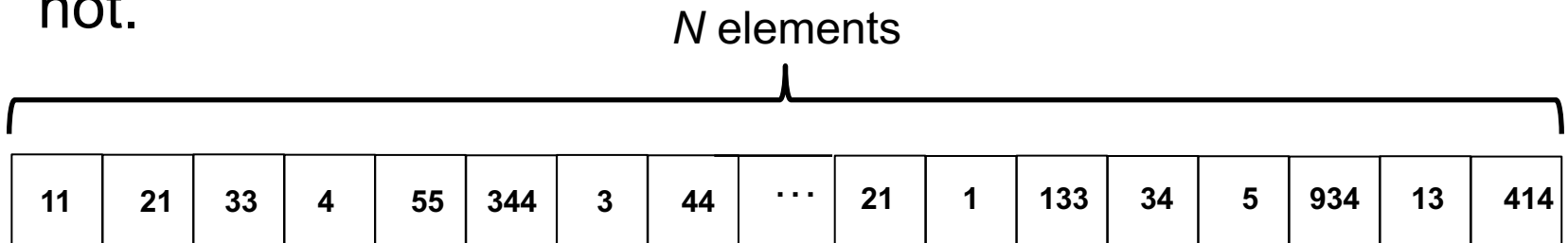
- An array of size n is filled with unsorted randomly placed, numbers.
- Goal: Find whether a specific item x is in this array or not.



↑
54?

Searching an Array of Size N for A specific Element

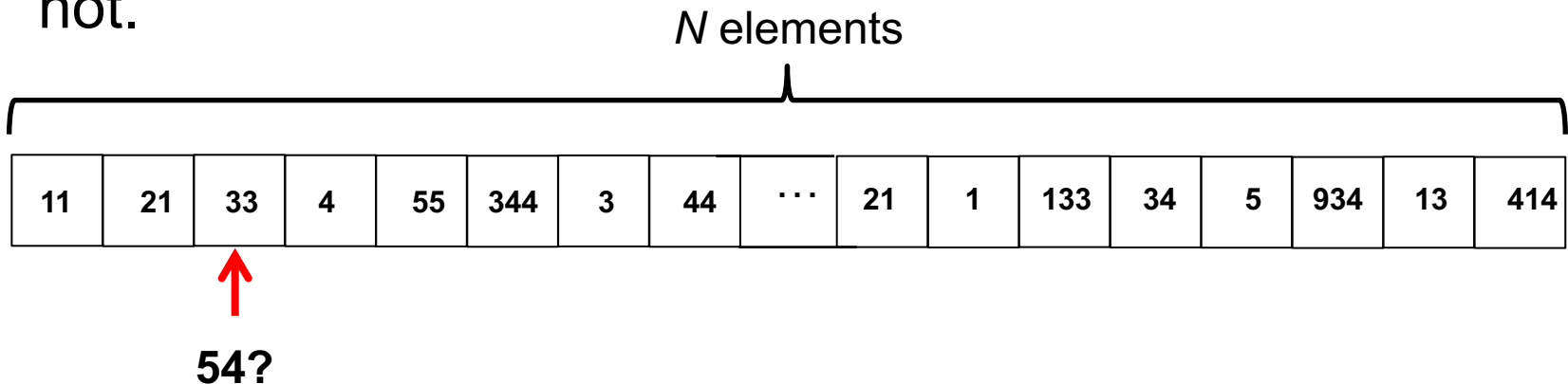
- An array of size n is filled with unsorted randomly placed, numbers.
- Goal: Find whether a specific item x is in this array or not.




54?

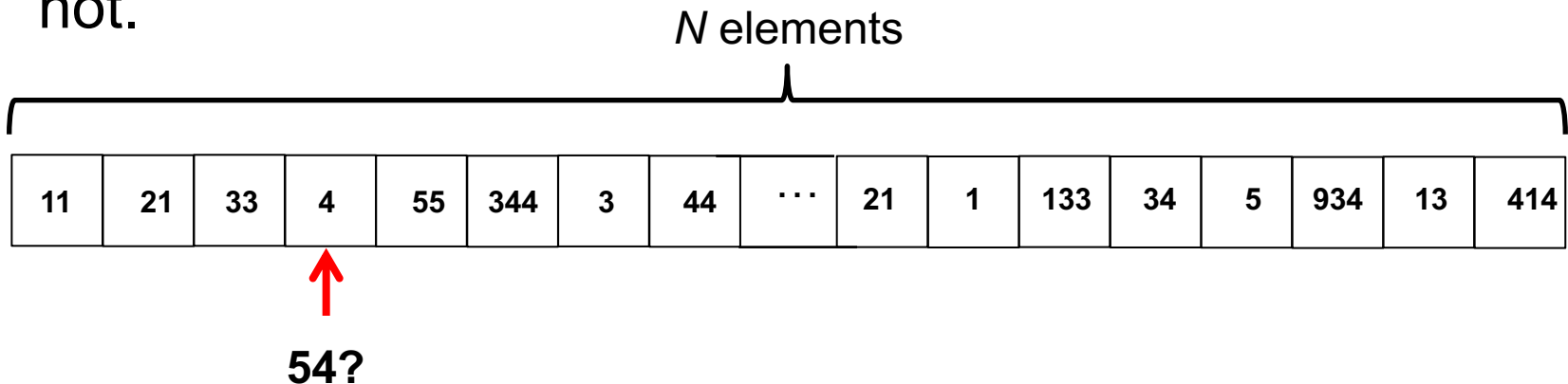
Searching an Array of Size N for A specific Element

- An array of size n is filled with unsorted randomly placed, numbers.
- Goal: Find whether a specific item x is in this array or not.



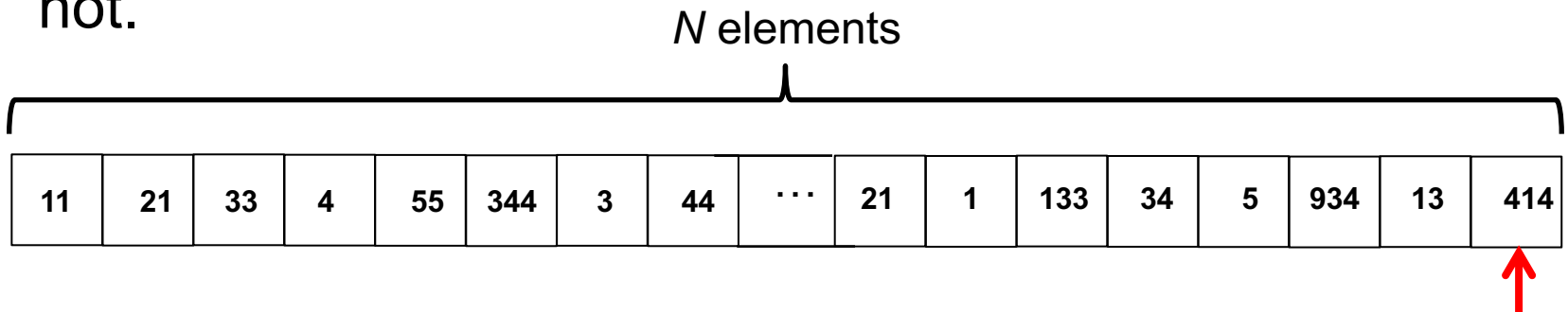
Searching an Array of Size N for A specific Element

- An array of size n is filled with unsorted randomly placed, numbers.
- Goal: Find whether a specific item x is in this array or not.



Searching an Array of Size N for A specific Element

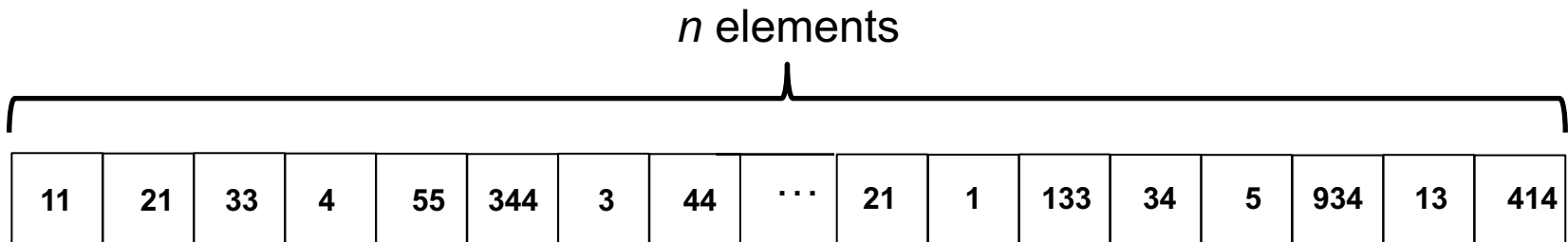
- An array of size n is filled with unsorted randomly placed, numbers.
- Goal: Find whether a specific item x is in this array or not.



54?

Searching an Array of Size N for A specific Element

- An array of size n is filled with unsorted randomly placed, numbers.
- Goal: Find whether a specific item x is in this array or not.



- We observe that this problem can be solved in n steps, where n is the size of the input.
- This problem belongs to computational complexity class of **P**.
- The class **P** is composed of such problems that can be solved in **P**olynomial time (or number of steps) with respect to the size of the input. ($O(n)=n, n^2, n^3, \dots$)

Complexity Classes

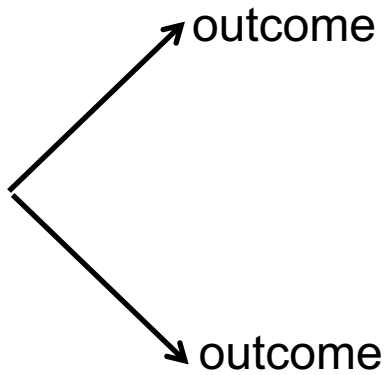
P: Can be solved in polynomial time.



A Series of Decisions to Obtain a Desired Outcome.

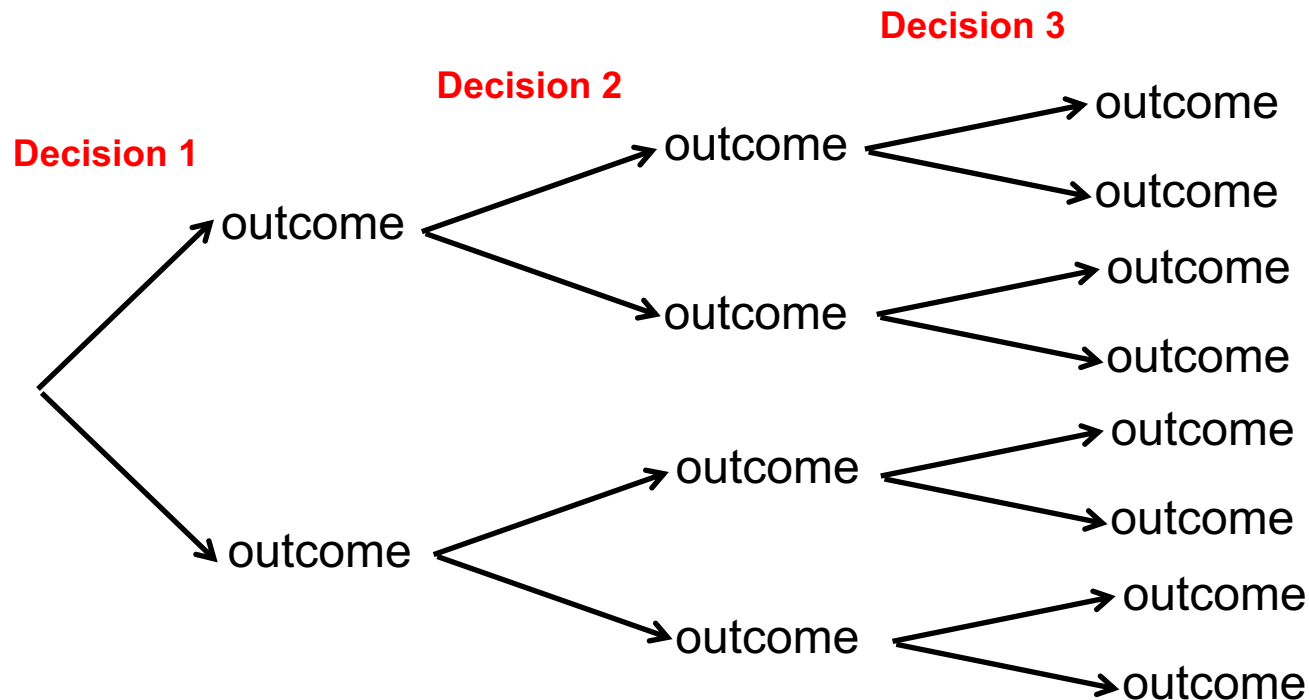
- Imagine making a decision with two possible outcomes.

Decision 1



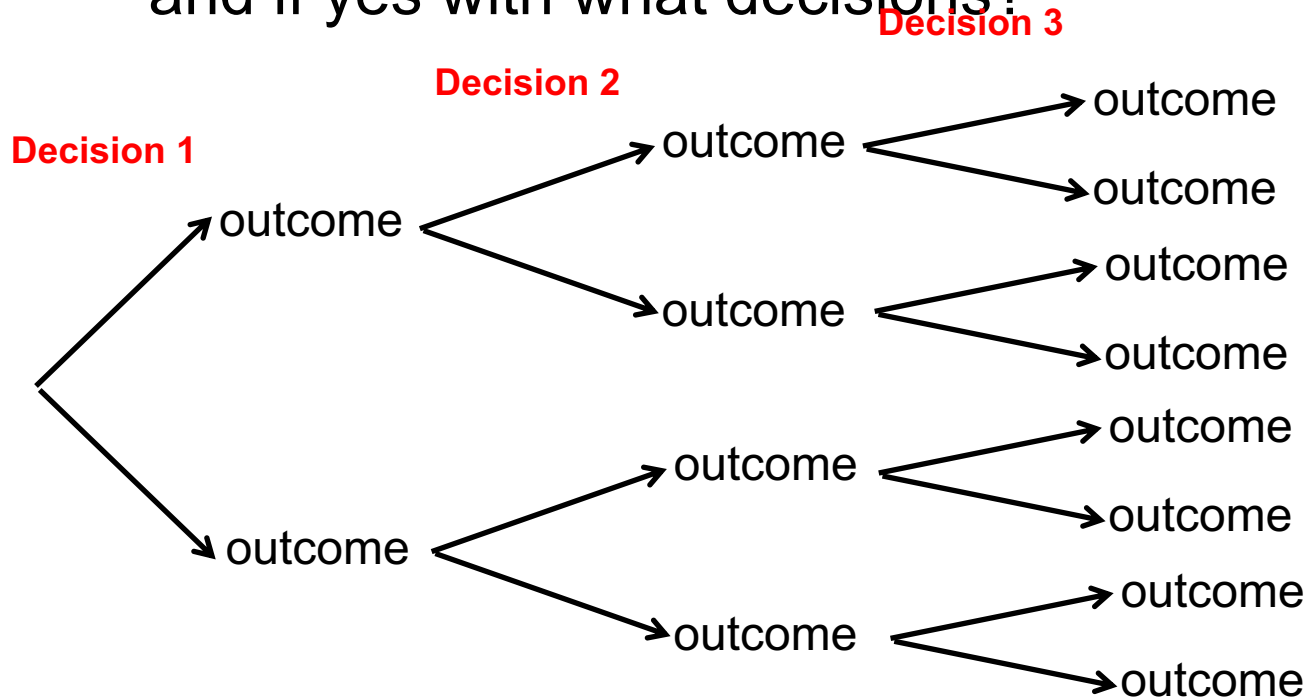
A Series of Decisions to Obtain a Desired Outcome.

- Imagine making a decision with two possible outcomes. And making further decisions with two outcomes for each decision.



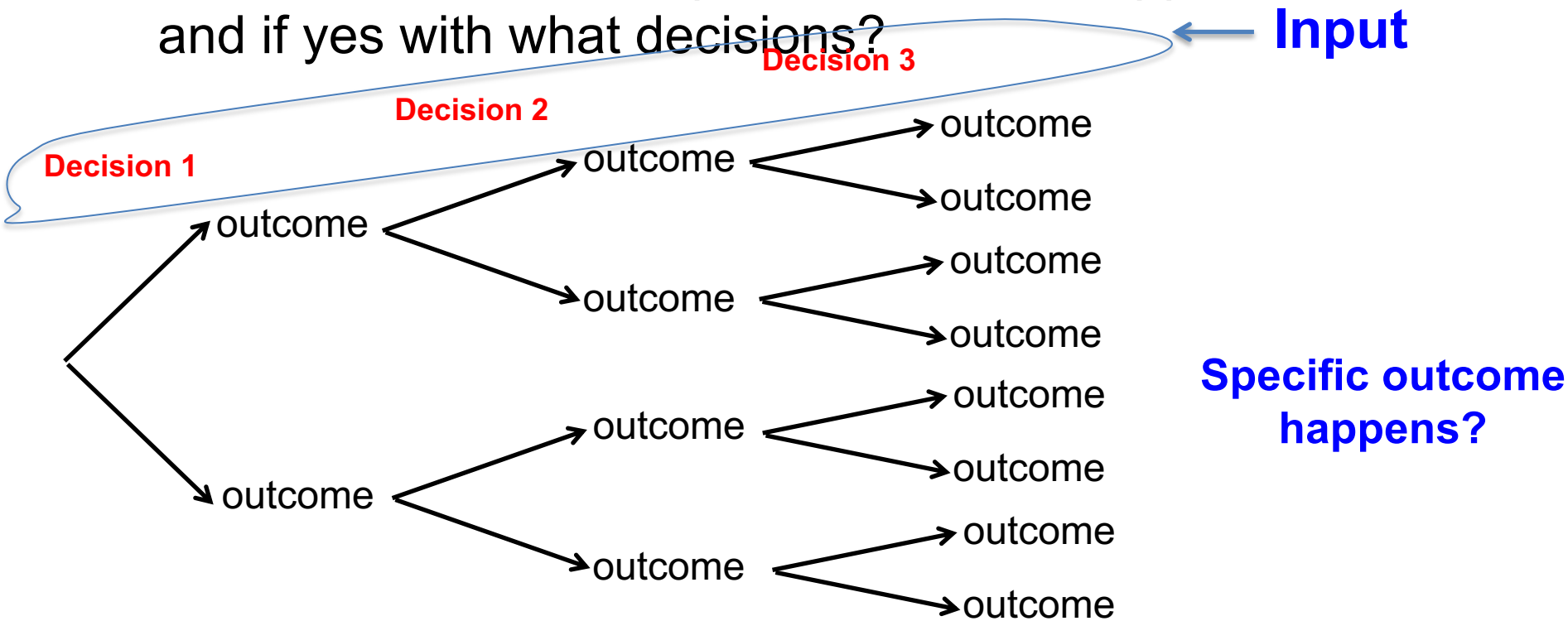
A Series of Decisions to Obtain a Desired Outcome.

- Imagine making a decision with two possible outcomes. And making further decisions with two outcomes for each decision.
- Goal: Find whether a specific outcome happens or not, and if yes with what decisions?



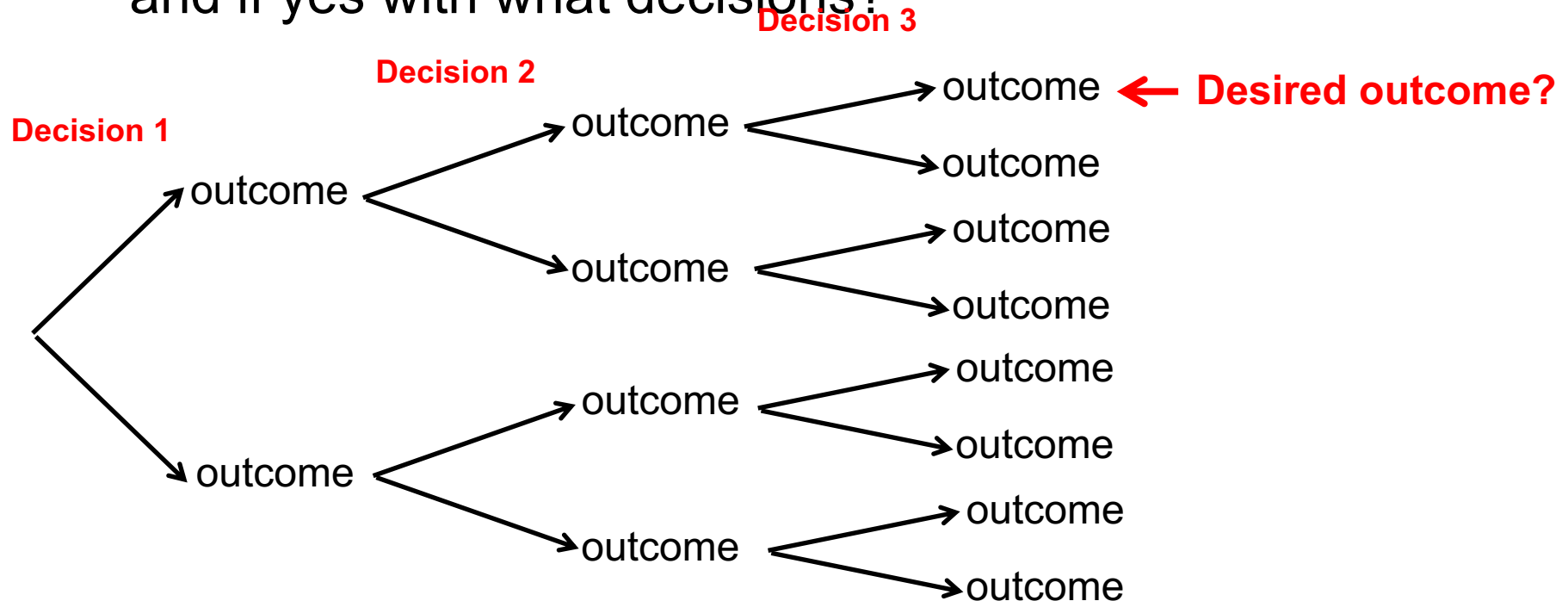
A Series of Decisions to Obtain a Desired Outcome.

- Imagine making a decision with two possible outcomes. And making further decisions with two outcomes for each decision.
- Goal: Find whether a specific outcome happens or not, and if yes with what decisions?



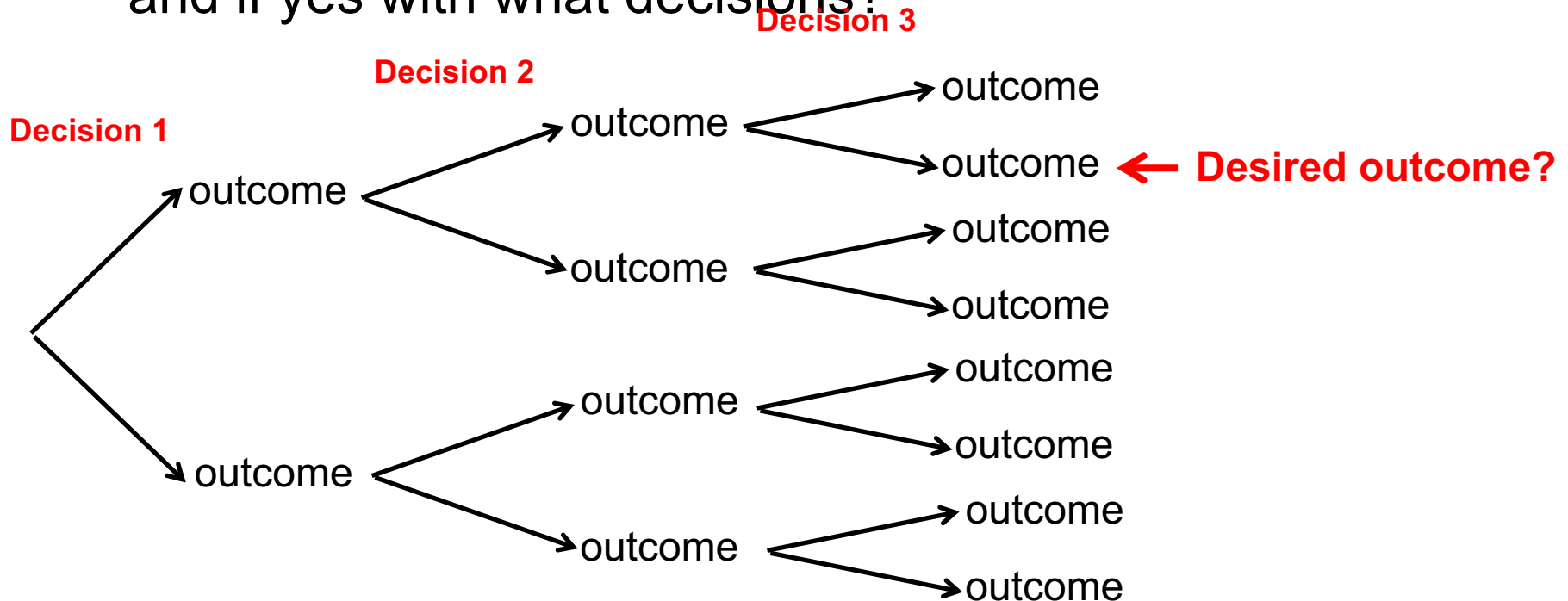
A Series of Decisions to Obtain a Desired Outcome.

- Imagine making a decision with two possible outcomes. And making further decisions with two outcomes for each decision.
- Goal: Find whether a specific outcome happens or not, and if yes with what decisions?



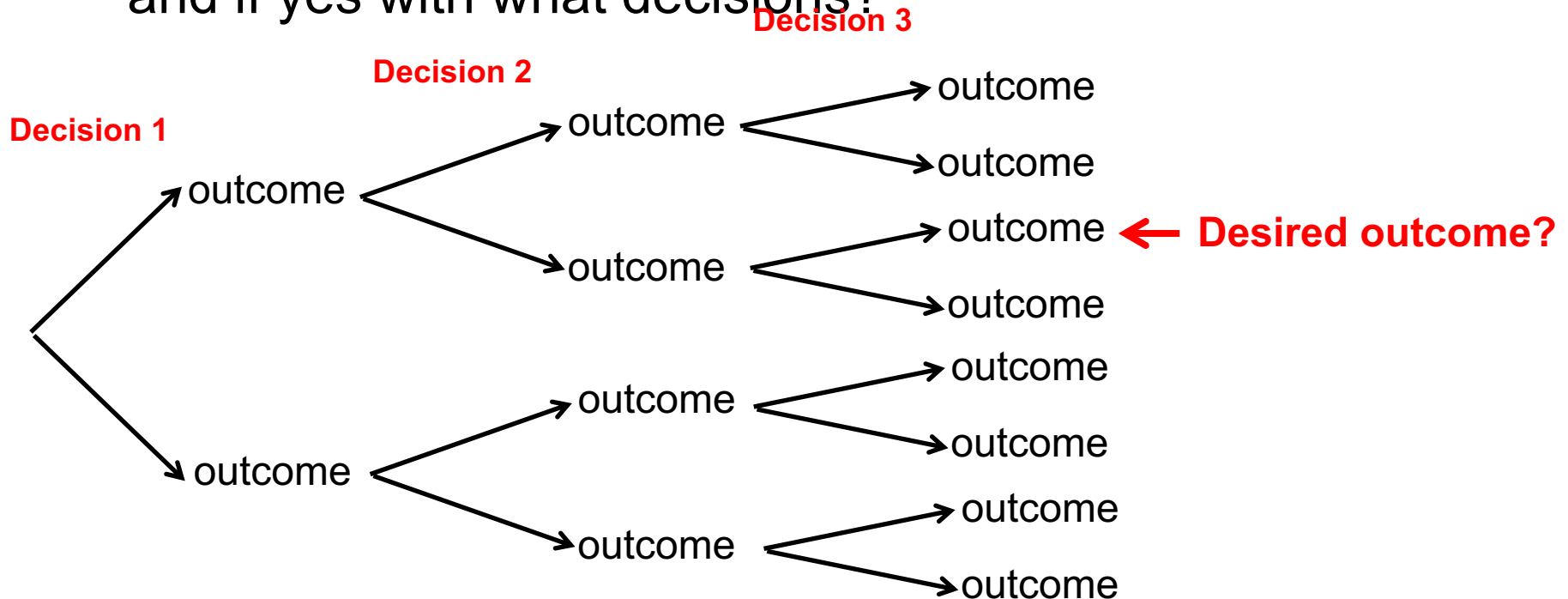
A Series of Decisions to Obtain a Desired Outcome.

- Imagine making a decision with two possible options. And making further decisions with two outcomes for each decision.
- Goal: Find whether a specific outcome happens or not, and if yes with what decisions?



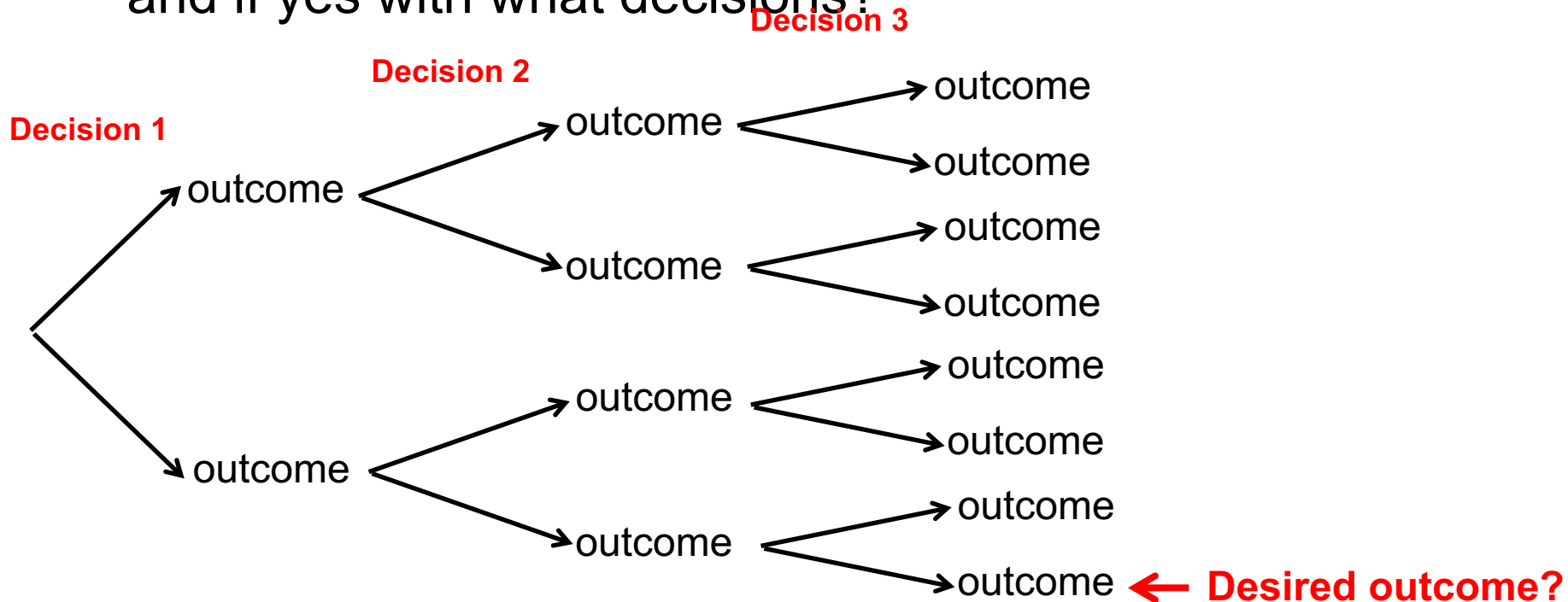
A Series of Decisions to Obtain a Desired Outcome.

- Imagine making a decision with two possible options. And making further decisions with two outcomes for each decision.
- Goal: Find whether a specific outcome happens or not, and if yes with what decisions?



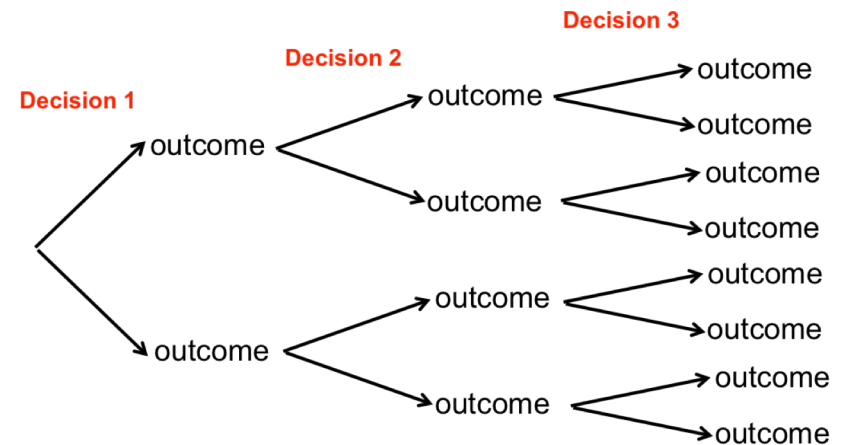
A Series of Decisions to Obtain a Desired Outcome.

- Imagine making a decision with two possible options. And making further decisions with two outcomes for each decision.
- Goal: Find whether a specific outcome happens or not, and if yes with what decisions?



A Series of Decisions to Obtain a Desired Outcome.

- Imagine making a decision with two possible outcomes. And making further decisions with two outcomes for each decision.
- Goal: Find whether a specific outcome happens or not, and if yes with what decisions?
- Here the number of decisions is the size of the problem.
- The number of possible outcomes to be checked exponentially increases by the size of the input.

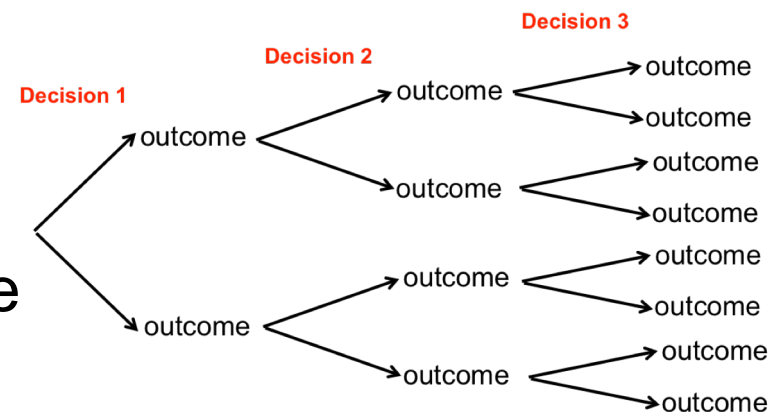


A Series of Decisions to Obtain a Desired Outcome.

- This problem belongs to complexity class of **NP**.
- **NP** is a class of decision problems that:
 1. A possible solution (in this case a series of decisions) can be verified in polynomial time.
 2. The problem can be solved in **N**on-deterministic **P**olynomial time.

Informal version of (2): We need a “lucky” algorithm to solve in polynomial time. In practice we cannot design a lucky algorithm.

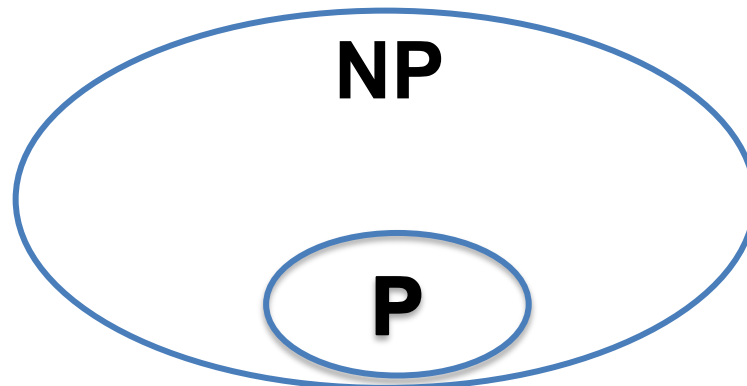
So far practical algorithms require exponential time.



Complexity Classes

P: Can be solved in polynomial time.

NP: Solutions verifiable in polynomial time. Can be solved in-deterministic polynomial time. (informal definition) No polynomial solution yet.

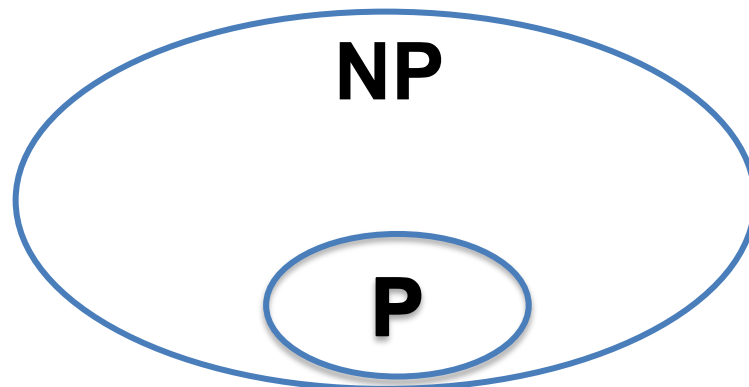


Complexity Classes

P: Can be solved in polynomial time.

NP: Solutions verifiable in polynomial time. Can be solved in-deterministic polynomial time. (informal definition) No polynomial solution yet.

Not all NP problems are equally hard!

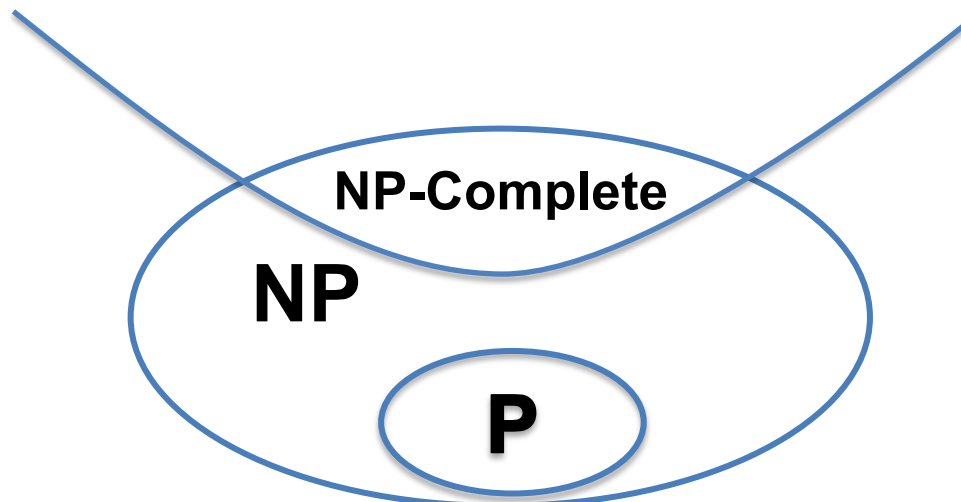


Complexity Classes

P: Can be solved in polynomial time.

NP: Solutions verifiable in polynomial time. Can be solved in-deterministic polynomial time. (informal definition) No polynomial solution yet.

NP-Complete: The hardest NP problems.



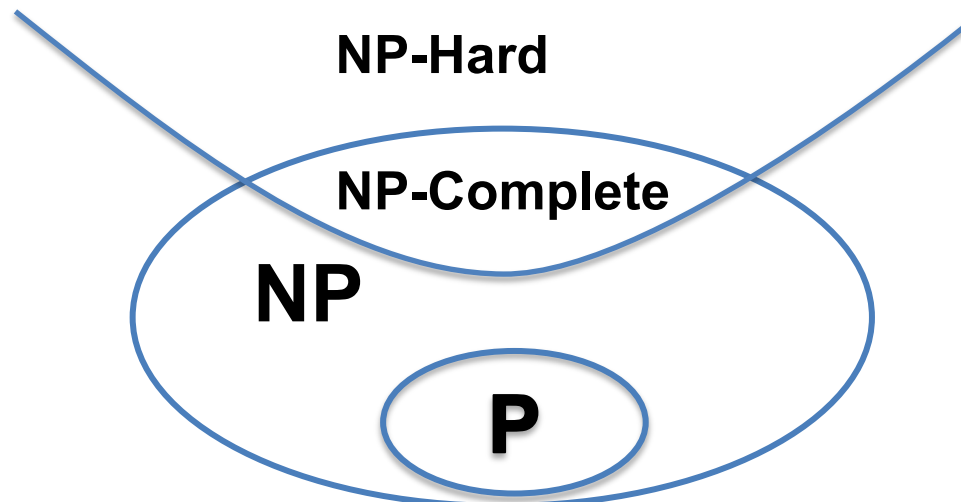
Complexity Classes

P: Can be solved in polynomial time.

NP: Solutions verifiable in polynomial time. Can be solved in-deterministic polynomial time. (informal definition) No polynomial solution yet.

NP-Complete: The hardest NP problems.

NP-Hard: At least as hard as the hardest NP problem.



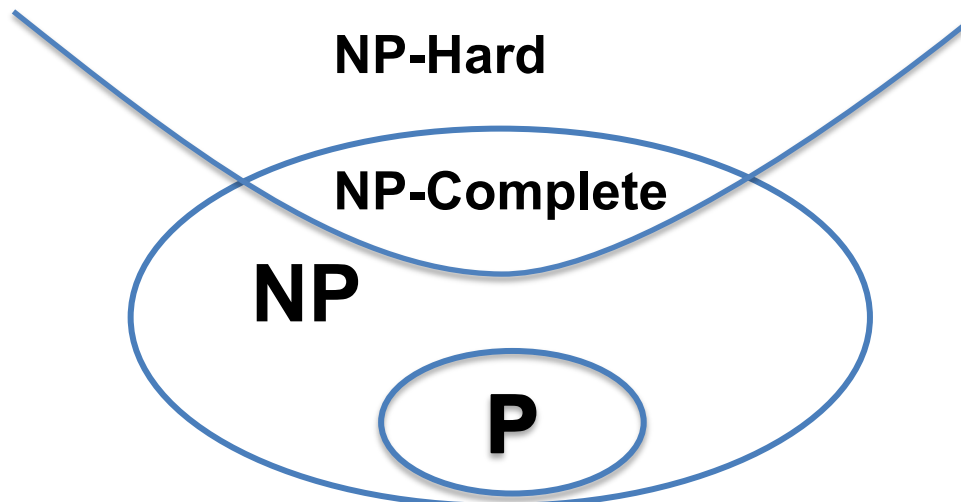
Complexity Classes: One Million Dollar Question!

P: Can be solved in polynomial time.

NP: Solutions verifiable in polynomial time. Can be solved in-deterministic polynomial time. (informal definition) No polynomial solution yet.

NP-Complete: The hardest NP problems.

NP-Hard: As hard as the hardest NP problem.



P=NP?

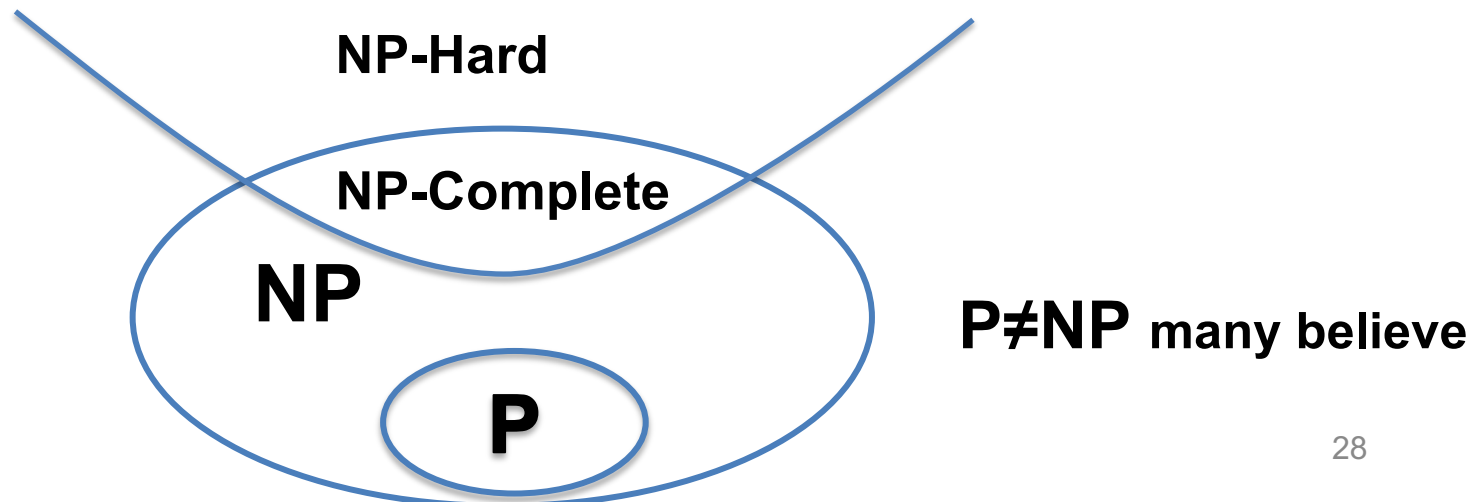
Complexity Classes: One Million Dollar Question!

P: Can be solved in polynomial time.

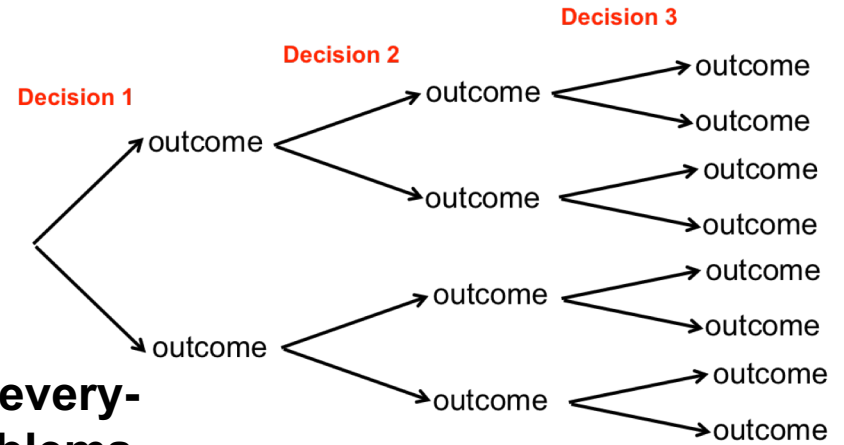
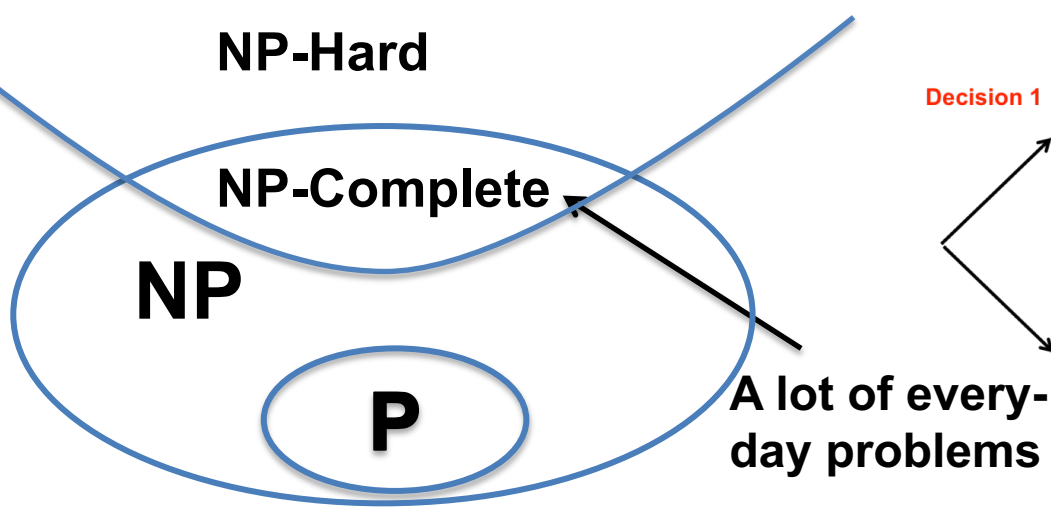
NP: Solutions verifiable in polynomial time. Can be solved in-deterministic polynomial time. (informal definition) No polynomial solution yet.

NP-Complete: The hardest NP problems.

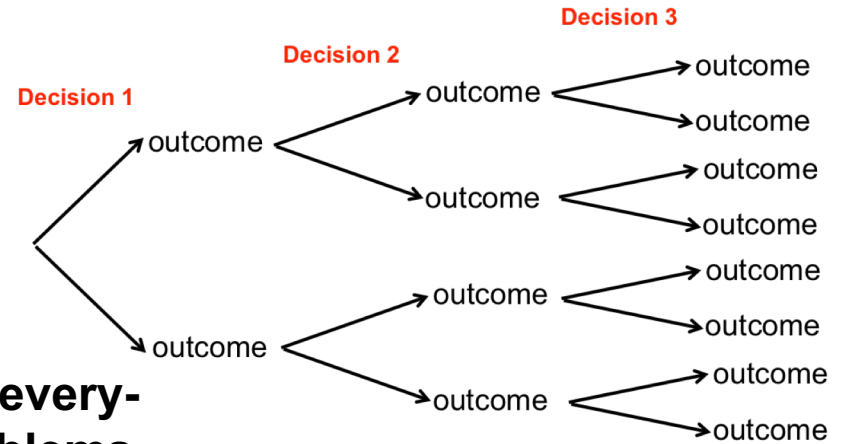
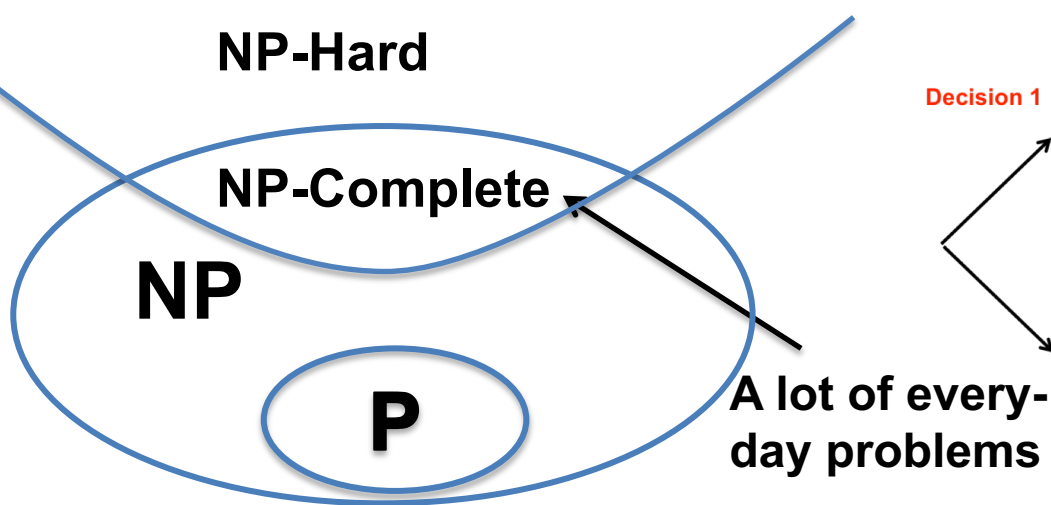
NP-Hard: As hard as the hardest NP problem.



$N \neq NP$ Implications

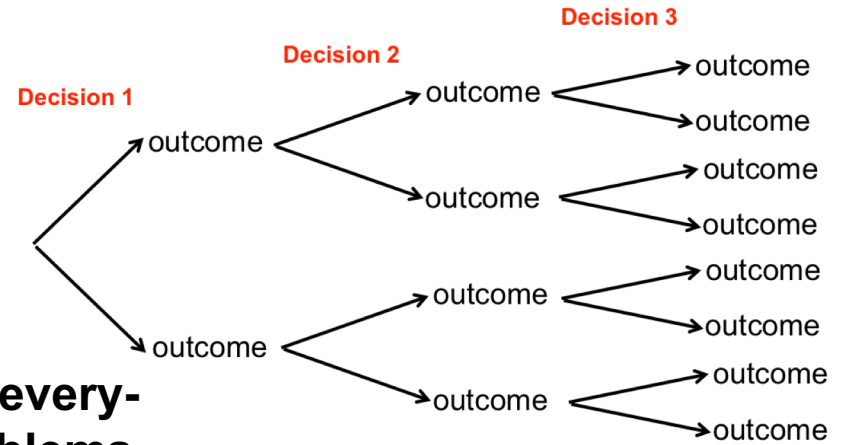
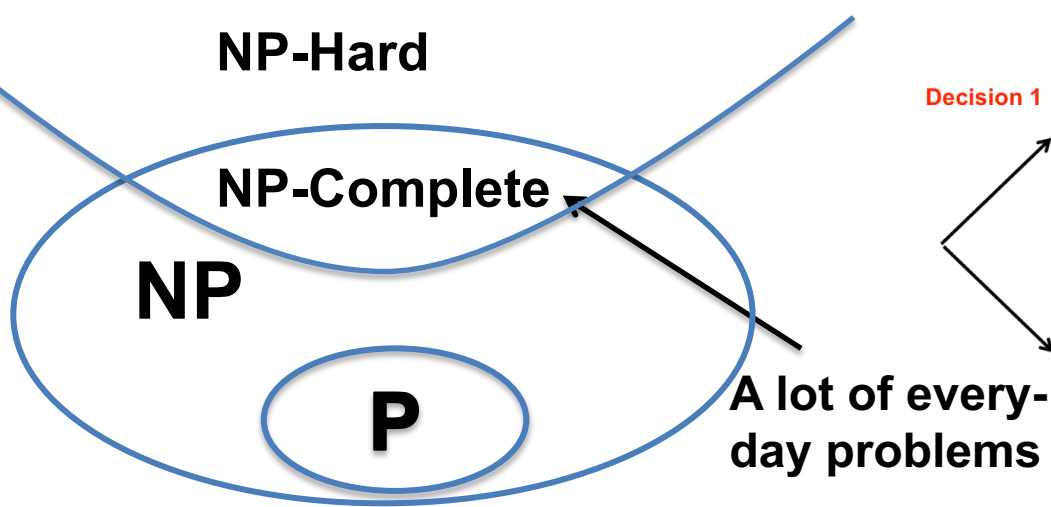


N≠NP Implications

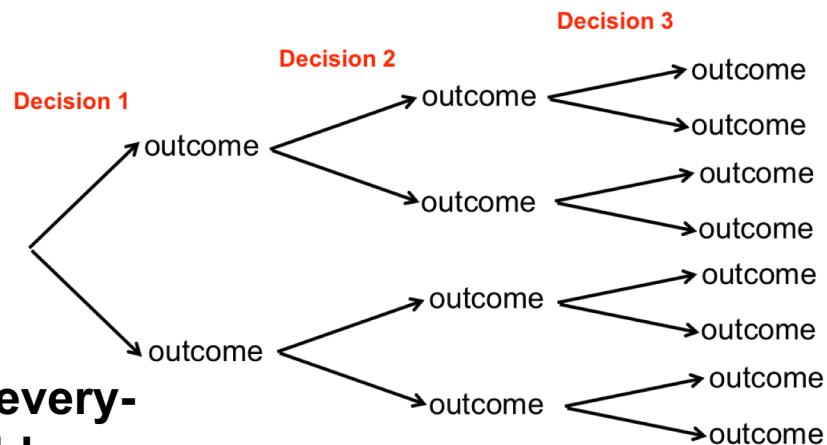
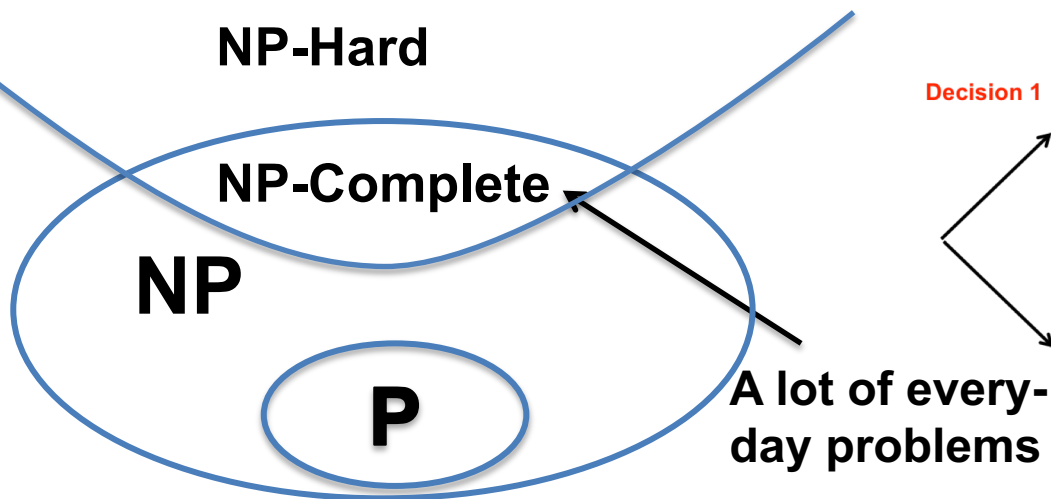


5	3			7			
6			1	9	5		
	9	8					6
8				6			3
4			8		3		1
7				2			6
	6					2	8
			4	1	9		5
				8			7
						7	9

Complexity Classes



Complexity Classes

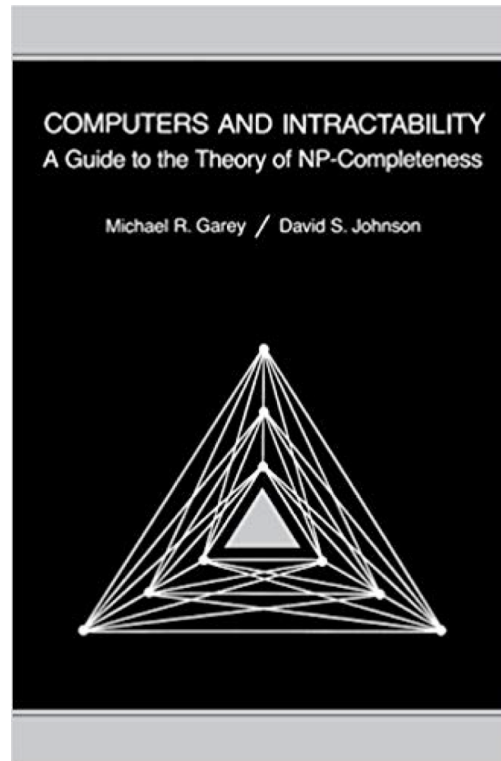


Strange: I went forward in time to see all the possible outcomes of the present situation.

Peter: How many did you see?

Strange: Fourteen million six hundred and five.

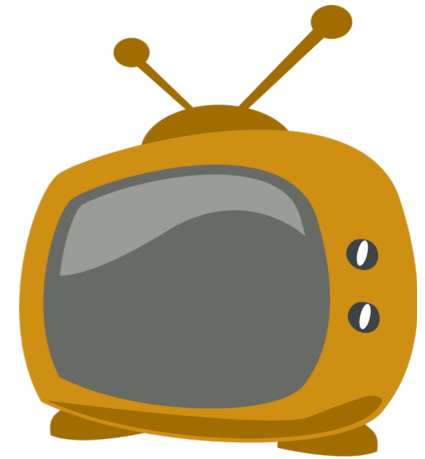
$$\text{Log}_2(14,000,605) \approx 24$$



Computers and intractability : a guide to the
theory of NP-completeness
Michael R. Garey, David S. Johnson.

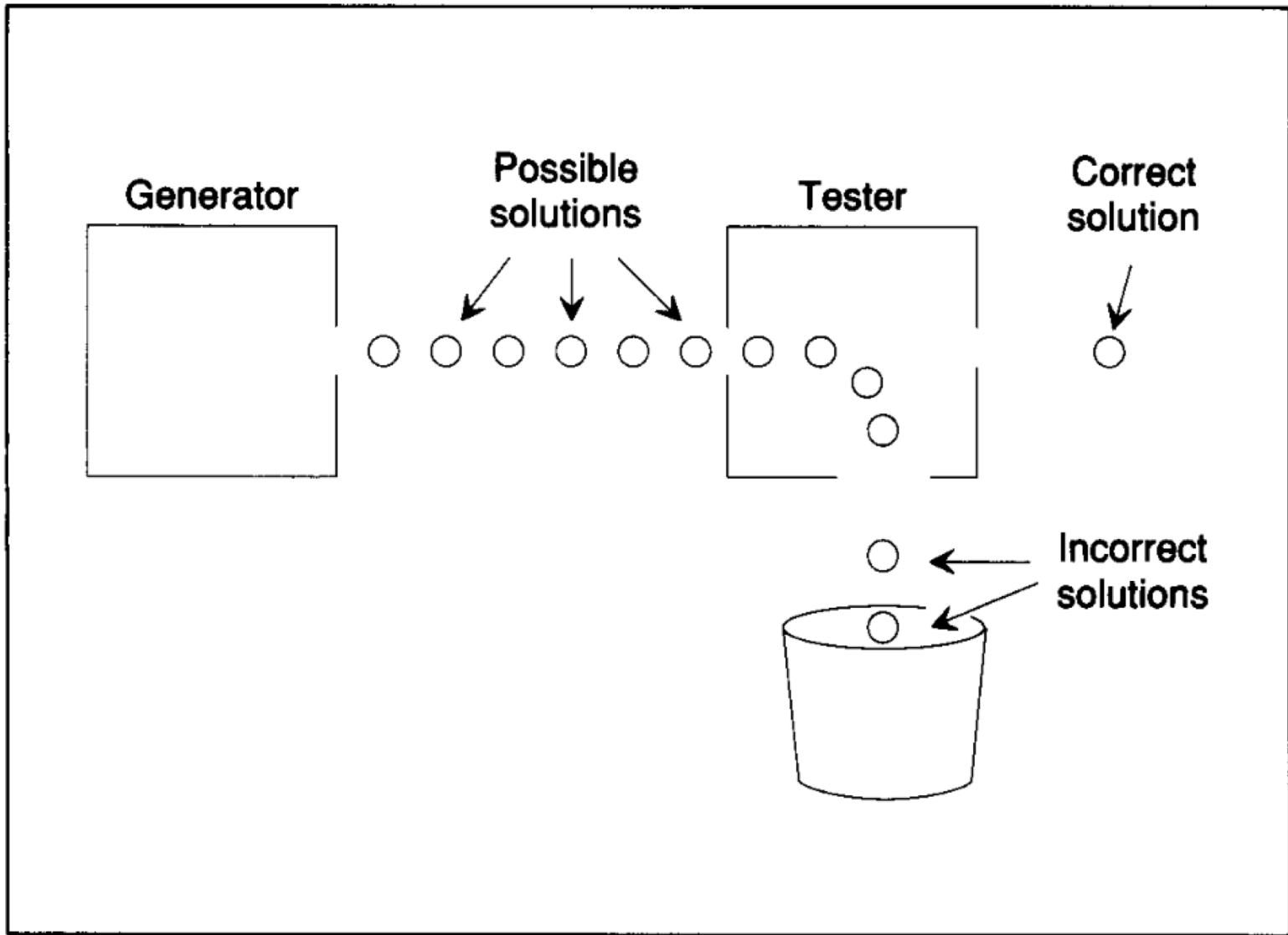
- Computational complexity is about the order of growth with respect to input size, not absolute time given a specific sized input.

video



$$\begin{aligned} \nabla \cdot \mathbf{E} &= \frac{\partial E_x}{\partial x} + \frac{\partial E_y}{\partial y} + \frac{\partial E_z}{\partial z} \\ (\nabla \times \mathbf{E})_x &= \frac{\partial E_z}{\partial y} - \frac{\partial E_y}{\partial z} \\ (\nabla \times \mathbf{E})_y &= \frac{\partial E_x}{\partial z} - \frac{\partial E_z}{\partial x} \\ (\nabla \times \mathbf{E})_z &= \frac{\partial E_y}{\partial x} - \frac{\partial E_x}{\partial y} \\ \nabla \cdot \mathbf{E} &= \frac{\rho}{\epsilon_0} & \nabla \times \mathbf{E} &= -\frac{\partial \mathbf{B}}{\partial t} \\ \nabla \cdot \mathbf{B} &= 0 & c^2 \nabla \times \mathbf{B} &= \frac{\mathbf{j}}{\epsilon_0} + \frac{\partial \mathbf{E}}{\partial t} \\ E &= \frac{q}{4\pi\epsilon_0} \left[\frac{\mathbf{e}_r'}{r'^2} + \frac{r'}{c} \frac{d}{dt} \left(\frac{\mathbf{e}_r'}{r'^2} \right) + \frac{1}{c^2} \frac{d^2}{dt^2} \mathbf{e}_r' \right], \\ c\mathbf{B} &= \mathbf{e}_r' \times \mathbf{E}. \end{aligned}$$

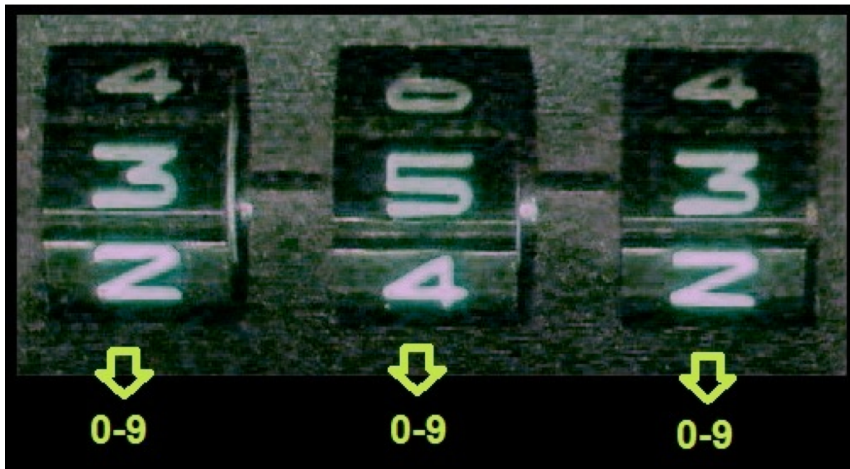
Search as a tool for problem solving and AI



Picture from Artificial Intelligence, Patrick Winston

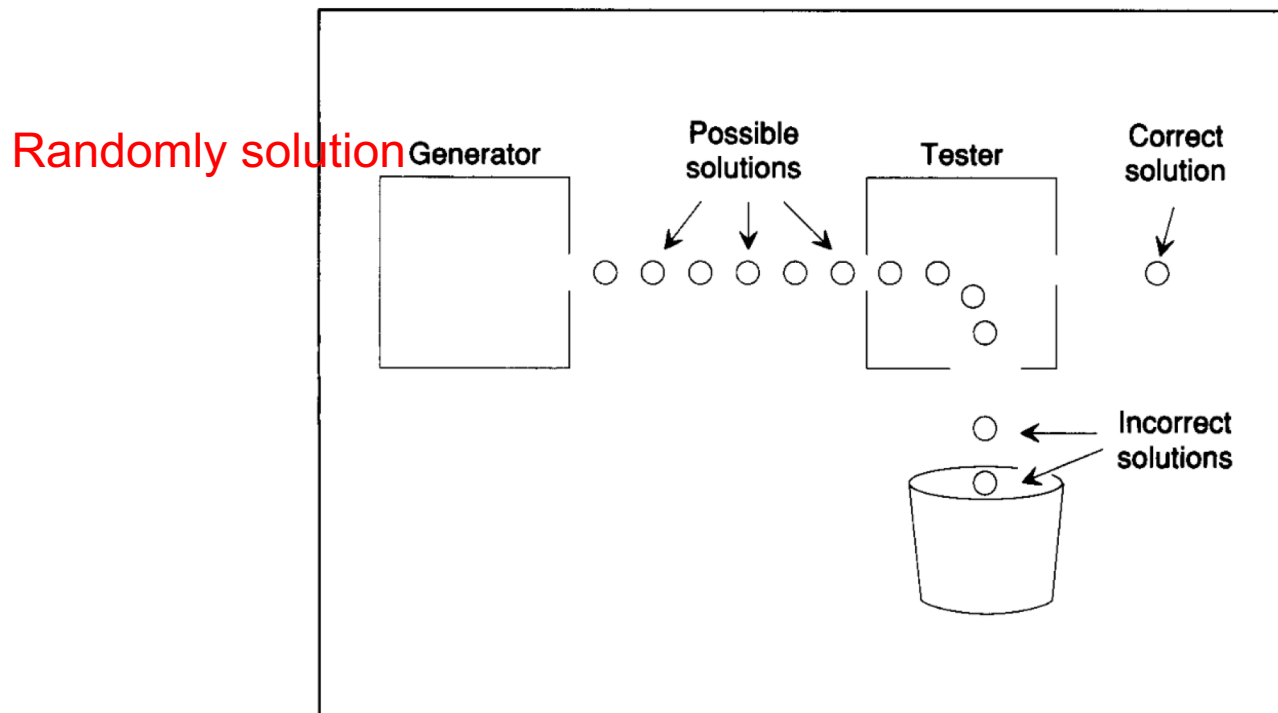
Search Mechanisms

Brute-Force (Exhaustive) Search



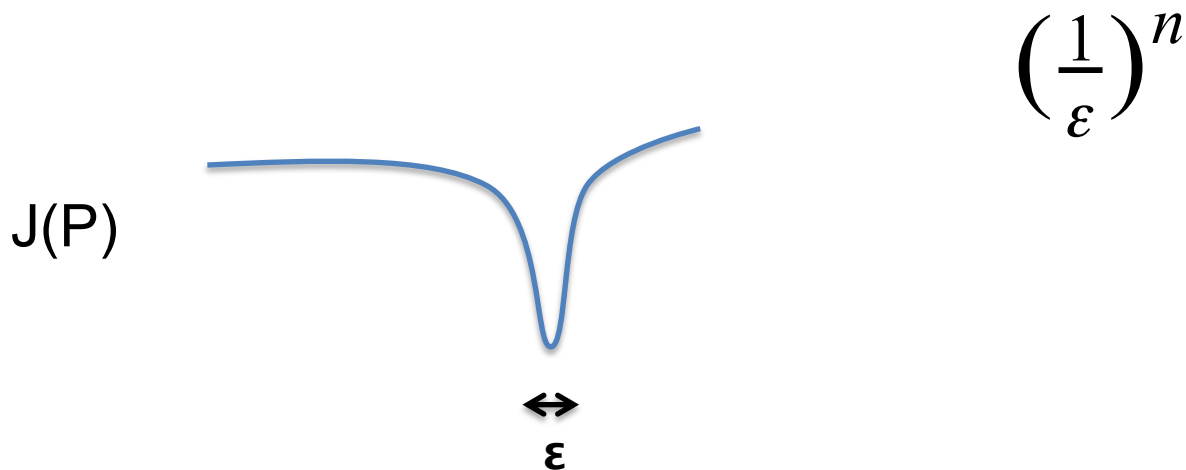
Random Search Methods: Pure Random Search

- Randomly (with a uniform distribution) choose candidates in the solution space up until max number of iterations performed, or an adequate fitness reached.



Random Search Methods: Pure Random Search

- Performance of pure random search method experiences large variation.
- In an n -dimensional function J , the expected number of iterations until a pure random sample point falls within the ε neighborhood of the minimal point of J is:



- The process of learning a function can be considered as a search in Euclidean space for a set of weights that implements the function.